# Call Admission Control Scheme to Determine the Preference for WLAN

Samrin Shareef, P.Prasanna Murali Krishna

Abstract— Call Admission Control (CAC) between heterogeneous networks plays an important role to utilize the system resources in a more efficient way. In this paper, we propose that the preference to the WLAN is determined based on the traffic load in the WLAN and the location of the cellular users. Our analysis relies on a previous study that divides the 3G cellular coverage area into zones based on the amount of resources that are required to support a connection to a mobile user. Using this model, we derive new call blocking and handoff failure probabilities as well as new call and handoff attempt failure probabilities.

Index Terms— CAC, vertical handoffs, heterogeneous networks, ring-based model, upward VHO, downward VHO, Call Blocking.

## 1 Introduction

ODERN wireless communication systems consist of Leellular, WiMAX, and WLAN access networks, coexisting in so-called heterogeneous networks. Multi-mode users can access multiple systems to get ubiquitous access in an efficient way in terms of cost and quality. From the network point of view, it is important to efficiently manage the resources, by considering vertical handoffs (VHOs) and call arrivals in the multiple-coverage area. Recently, the 3rd Generation Partnership Project (3GPP) Long Term Evolution (LTE) system has gained lots of interest for future broadband wireless networks. Consequently, its integration with Wireless Local Area Networks (WLANs), including seamless/vertical handoffs and Call Admission Control (CAC), has held the spotlight in the heterogeneous networks. Unlike WLANs, where broadband data transmission is available at low cost and features a simple control plane, 3GPP LTE systems can implement more complex resource management schemes to provide more efficient services and better Quality of Service (QoS) with a typical cell radius of less than 300 m [1].

There have been some works in cellular/WLAN integrated networks [4], [5]. A simple admission strategy is to have an unconditional preference to WLANs since WLANs are cheaper and have more bandwidth compared to the early version of 3G networks [4]. However, in 3GPP LTE systems, the bandwidth and cost have become comparable to that of WLANs due to wider bandwidth adoption and flat-rate tariffs. Moreover, WLANs may be over-crowded due to the WLAN first scheme, and thus a large portion of connection attempts to WLANs can fail.

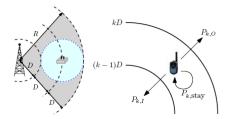
In [5], the authors considered that users with high mobility tend to spend a short time in the double-coverage area and require an *upward VHO* back to the cellular network shortly after a *downward VHO* to the WLAN. This may cause unnecessary VHO processing load in both networks. Therefore, by probabilistically rejecting VHOs for highly mobile cellular

users, the VHO processing load and VHO blocking probability can be reduced while maintaining reasonable throughput in the WLAN.

We proposed a ring-based cell model in [2], where the required resources depend on which annular region of the cell the user occupies. Unlike conventional Frequency/Time/Code Division Multiple Access. Networks, more resources are allocated to a user farther from the base station (BS) in terms of power, bandwidth, and time slots. We found that the ring-based model allows more accurate analysis of the system performance, especially with respect to handoff and mobility metrics.

It may not be beneficial for the cellular users in the dual coverage area to have a policy of choosing the WLAN with a fixed probability regardless of the load on the WLAN or the location of the user. For example, it is desirable that a user located farther from the BS access the WLAN since more resources would be released to the cellular network when the user migrates to the WLAN. In this paper, we develop a traffic management policy which sets the probability of migrating to the WLAN based on the traffic load in the WLAN and the location of the user in the cellular system. Using the ring-based model with a single WLAN hotspot, we compute WLAN migration preference values as functions of the blocking and handoff failure probabilities as well as new call/handoff attempt failure probabilities. The network can perform efficient call admission control by adaptively changing the WLAN preference settings in heterogeneous wireless networks in response to changing traffic conditions.

The rest of this paper is organized as follows. In Section II, we describe the system model, where mobility-related parameters are obtained using a ring-based model and system performance is derived based on steady-state analysis. System performances are investigated through intensive simulations in Section III. Finally, we present our conclusions in Section IV.



(a) A sector divided (b) Transition probabilities in the *k*th ring into rings, *K*=3, with a WLAN hotspot

Fig. 1. System model.

## 2 System Modeling

A cell has three sectors and each sector is divided into K concentric rings, each of whose width is D m. Fig. 1(a) shows an example sector with 3 rings, where a single WLAN AP, whose coverage radius is D m, is located 2D m from the BS. In the double coverage area, a cellular user can access either the cellular network or a WLAN [4], and can perform handoffs to move from one to the other. These vertical handoffs (VHOs) can be categorized into downward VHOs and upward VHOs [5]. The downward VHO is a handoff from the cellular network to the WLAN; the upward VHO is a handoff in the opposite direction.

# 2.1 Modified Ring-Based Model

Let  $n_k$  be the number of mobile users in the kth ring in the cellular network, for  $k=1,\ldots,n$   $n_{k+1}$  is the number of cellular users being serviced in the WLAN. The system state S is denoted as  $(n1,\ldots,nK,nK+1)$ . Ac is the area covered by a 3G sector, and Ac,k is the area of the kth ring of a sector and Aw,k is the area of the WLAN that overlaps the kth ring,  $k=1,\ldots,K$ . We can define Pc, and Pw, as the probabilities that a user lies in the kth ring excluding the region of overlap with the WLAN and in the intersection of the kth ring and the WLAN, respectively. Assuming a uniform mobility, these probabilities can be written as

$$P(w,k) = \frac{A(w,k)}{Ac}$$
 (1)

$$P(c,k) = \frac{(A(c,k) - A(w,k))}{Ac}$$
(2)

for  $k = 1, \dots, K$ ,  $P_{n\sigma} = \sum_{k=1}^{K} P_{n\sigma,k}$  and  $P_{n\sigma} = \sum_{k}^{K} P_{n\sigma,k}$  and therefore,  $P_{n\sigma} + P_{n\sigma} = 1$ .

Let  $\lambda c$  be the new call arrival rate in a sector for the cellular network and  $\lambda h$  be the handoff rate of users into the sector. and the average call holding time for a user is denoted as  $1/\mu c$ . Since WLAN users tend to remain stationary, it is assumed that there is no handoff request for WLAN users. In the WLAN, all users are assumed to require the same amount of resources for simplicity [5]. For the cellular network, however, a user's required resources vary

depending on its distance from its BS, and guard channels are reserved for handoff and mobility [2]. In this paper, we consider CBR traffic for simplicity but extending the analysis to multiple traffic types is straightforward. When a call arrives in the kth ring in the cellular network, it requires rk resources, while a user in the WLAN requires rw resources. Cc and Cw are the capacities of the cellular network and the WLAN, respectively, and Cg is the amount of guard resources reserved by the cellular network. A fixed amount of resource, nw is allocated to the WLAN users and the remaining resource of Cw - nw can be assigned for the cellular users who request connections to the WLAN.

Rc(S) and Rw(S) are the total resources being used in the cellular network and the WLAN, respectively, when the system is in state S. They are denoted by

$$Rc(S) = \sum_{k=1}^{K} nkrk$$
 (3)

$$Rw(S) = (n(K+1) + nw) rw$$
 (4)

# 2.2 Mobility-Related Parameters

First, we consider the mobility of users between neighboring rings. When a mobile enters a given ring, it is assumed to reside there for an exponentially distributed residual time whose mean is  $1/\gamma$ . At the end of that time period, a user in the kth ring may move to the (k-1)th or (k+1)th rings with probabilities of Pk, I or Pk, O, respectively (see Fig. 1(b)). Or the user may stay within the kth ring with probability Pk, stay, which can be equivalently seen as the user moves within the ring after a certain residual time. Assuming a user's direction of movement is uniformly distributed, we can obtain that Pk, stay = Pk, I+Pk, O=1/2. Pk, o and Pk, I are determined in terms of the inner and outer circumferences of the kth ring, Ik, I and Ik, I. For I0. For I1, the parameters are given by

$$P(k,I) = \frac{1}{2} l(k,I)/(l(k,I) + l(k,o))$$
 (5)

$$P(k, O) = \frac{1}{2} - P(k, I)$$
 (6)

These results agree with those in [2], where the probabilities  $\{P_{k+1}\}_{k=1}^n$  and  $\{P_{k+1}\}_{k=1}^n$  have been calculated by using the balance equations to satisfy the condition that the average inward-rate should be equal to the average outward-rate at each ring boundary. Exceptionally for the innermost "ring" (i.e., k=1), a user always moves out of the ring after a certain residual time, so P1,0 is always 1 and other probabilities become 0.

Now we consider the probabilities of moving between the cellular and WLAN areas. When a user moves out of the WLAN's coverage, it may move to the kth ring in the sector or it may move to neighboring cells. The corresponding probabilities are denoted as  $\beta k$  and  $\beta K+1$ . Assuming the mobile's direction is uniform over  $[0, 2\pi]$ , these probabilities are respectively proportional to the length of the WLAN's boundary that lies within the kth ring or neighboring cell. Generally, we require a degree of margin in the received signal strength (RSS) for handoff to prevent ping-pong

effects. Assuming a margin of  $\xi th$  dB,  $\beta 1$  and  $\beta K+1$  can be calculated assuming a given path loss. That is, for the first ring and the Kth ring, if a user leaves the WLAN and the path loss difference from the WLAN boundary to the corresponding sector/ring boundary is less than the margin, then the user moves to the first ring or to a neighboring cell. For example, we consider a simple 3-ring case with D=100 m, where the path loss (in dB) is given by 39.95+43.375 log10 d10 (din meters) and  $\xi th=1$  dB. Then,  $\beta 1=0.11$ ,  $\beta 2=0.31$ ,  $\beta 3=0.37$ , and  $\beta 4=0.21$ .

We next consider the balance equations between the cellular network and the WLAN. The users in the kth ring moves into the WLAN with a probability of  $\alpha k$ . On the boundary of the WLAN coverage within each of the K rings, the average move-in rate from the cellular network to the WLAN should be equal to the average move-out rate from the WLAN to the cellular network, i.e.,

$$Pc1\gamma 1\alpha 1 = Pw\gamma w\beta 1 \tag{7}$$

$$Pc2\gamma 2\alpha 2 = Pw\gamma w\beta 2 \tag{8}$$

.

 $Pc(n-1)\gamma(n-1)\alpha(n-1) = Pw\gamma w\beta(n-1) (9)$ 

$$\gamma ra(k+1) = \text{Pw}\gamma w\beta(k+1)$$
 (10)

where  $1/\gamma w$  and  $1/\gamma 1$  are the mean residual times in the WLAN and in the innermost ring, respectively3.  $\alpha K+1$  is the probability that a handoff call from a neighboring cell requests a VHO to the WLAN. The residual time is assumed to be proportional to the size of the region [2], [3], i.e.,  $1/\gamma R = \alpha R$ ,  $1/\gamma 1 = \alpha D$ , and  $1/\gamma = \alpha D 2$ , where  $1/\gamma R$  is the mean residual time in a sector, R is the cell radius, and  $\alpha$  is a constant determined by the mobility characteristics of the users.

## 2.3 Call Admission Control Scheme

In the WLAN, a fixed amount of resource, nw is allocated for WLAN users and the WLAN's remaining resource, Cw - nw can be assigned to cellular users that newly appear in the double-coverage area or that request downward VHOs. For the cellular network, however, the resources that are used by a mobile depend on the mobile's location, and guard resources are reserved to support handoff and mobility [2].

When a new call appears in the intersection of the WLAN coverage area and the kth ring, it is serviced by the cellular network with a probability of 1 - wk. With probability wk, a new call attempts to connect to the WLAN. If no resources are available in the WLAN, the mobile tries to connect to the cellular network; the call is blocked only when resources are not available in the cellular network and the WLAN.

For downward VHO requests by mobiles in the kth ring, they are processed in the WLAN with probability wk. The call continues to be serviced in the cellular network without requesting a VHO with probability 1 - wk. When the call is

successfully handed off to the WLAN, rk resources are released by the cellular network and rw resources are occupied in the WLAN. If  $\Delta r = rk - rw > 0$ , it is beneficial to initiate VHOs to reduce the overall resource usage. Larger  $\Delta r$  values indicate more efficient resource usage due to VHOs. Users from neighboring cells enter the double-coverage area with a probability of  $\alpha K+1$  and request downward VHOs to the WLAN with a probability of wK. If no resources are available in the WLAN, their requests are rejected and they try to access the cellular network.

We measure the system performance using the call blocking probability PB and handoff failure probability Pf [2]. The new call may be blocked in the cellular network or the WLAN with probabilities PB,C and PB,W, respectively, and PB = PB,c+PB,w. Similarly, Pf = Pf,c+Pf,w, where Pf,c and Pf,ware the probabilities of handover failure to the cell and WLAN, respectively. The goal of the CAC is to find the optimum value for the vector of WLAN choice probabilities, which we call the WLAN preference settings,  $W = [w1, \ldots,$ wK], so that PB and Pf are minimized. We introduce two new probabilities, P'B and P'f, as cost measures. P'B is the new call attempt failure probability, and P'f is the handoff attempt failure probability. When a new call or VHO call attempts to connect to the WLAN, its attempt fails if no resources are available in the WLAN. However, it is not blocked or lost unless there are no resources available in the cellular network as well. These attempt failures result in wasted resources in the control plane; thus it is also important to reduce these failure probabilities to minimize cost and complexity.

# 2.4 Steady-State Analysis

The state transitions are investigated as follows;

- New call arrivals outside the WLAN: When a call arrives in the non-overlapped kth ring at a rate of λc,k = λcPc,, nk → nk +1 if Rc(S) + rk ≤ Cc − Cg. Otherwise, it is blocked.
- New call arrivals in the double-coverage area: When a new call appears in the kth ring at a rate of  $\lambda cPw$ ,, it attempts to connect to the WLAN or the cellular network with probabilities of wk or 1 wk, respectively. When the new call uses the WLAN,  $nK+1 \rightarrow nK+1 + 1$  if  $(S)+rw \leq Cw$ . Otherwise,  $nk \rightarrow nk+1$  if  $(S)+rk \leq Cc Cg$ .

When a call terminates, its resources are released to the cellular network or the WLAN as follows.

- When a call terminates in the kth ring with a rate of  $nk\mu c$ ,  $nk \to nk 1$  for  $k = 1, \ldots, K$ .
- When a call terminates in the WLAN with a rate of  $nK+1\mu c$ ,  $nK+1 \rightarrow nK+1-1$ .

Handoffs from adjacent cells can be categorized into two cases. A handoff to a neighboring cell that arrives outside that cell's embedded WLAN occurs with a probability of  $1-\alpha K+1$ ; a handoff to a neighboring cell's double-coverage area occurs with a probability of  $\alpha K+1$ . The mobile that is handing off to the double-coverage area requests a VHO with probability Handoffs from adjacent cells can be categorized into two cases. A handoff to a neighboring cell that arrives outside that cell's embedded WLAN occurs with

a probability of  $1-\alpha K+1$ ; a handoff to a neighboring cell's double-coverage area occurs with a probability of  $\alpha K+1$ . The mobile that is handing off to the double-coverage area requests a VHO with probability

- Handoff calls move into the cellular-only area at a rate of (1 − αK+1)h; when a handoff of this type happens, nK → nK + 1 if Rc(S) + rK ≤ Cc.
- Calls perform VHOs to the WLAN in the double coverage area with a rate of  $\omega K\alpha K+1\lambda h$ ; when this happens,  $nK+1\to nK+1+1$  if  $Rw(S)+rw\leq Cw$ , or  $nK\to nK+1$  if  $Rc(S)+rK\leq Cc$ .
- Calls perform horizontal handoffs to the cellular network from the double-coverage area at a rate of (1 − ωK+1)h; when one of these occurs, nK → nK + 1 if Rc(S) + rK ≤ Cc.

During a downward VHO, a call in the kth ring attempts a VHO to the WLAN with a probability of wk. We must have  $Rw(S)+rw \le Cw$  in the WLAN, after accepting the request. After the VHO completes, rk resources are released in the cellular network. Therefore, the larger  $\Delta r$  is, the greater wk should be. For upward VHOs, rw resources are released in the WLAN and rk resources are occupied in the sector.

- Calls in the kth ring (k = 2, ....) moves to the double coverage area at a rate of nkαkwkγ and change S as follows: nk → nk-1 and nw → nw+1, if Rw(S)+rw ≤ Cw.
- Calls in the first ring move to the double-coverage area at a rate of n1α1w1γ1, changes in S are the same as above.

Calls move out to the kth ring at a rate of  $nK+1\gamma w\beta k$ , k=1, . .,, and S changes as follows:  $nK+1 \rightarrow nK+1-1$  and  $nk \rightarrow nk+1$ , if  $Rc(S)+rk \leq Cc$ .

Calls are handed off to neighboring cells at a rate of  $nK+1\gamma w\beta K+1$ , and S changes as follows: nK+1-1.

When a user moves outward or inward between rings in the cellular network, it requires more resources or fewer, respectively [2].

- When a user in the kth ring (k = 1, . . .,K 1) moves outward, it requires more resources and the state Schanges as nk → nk-1 and nk+1 → nk+1+1 with a rate of n1γ1 for the first ring and nkPk,Oγ for the other rings, respectively. If (S) + δrk ≥ Cc, the user can maintain its connection, where δrk = rk+1 rk is the additional resources required by the user. In the outmost ring (theKth ring), the user can be handed off to a neighbor cell, i.e., nK → nK 1 with a rate of nKPK.
- For a ring (k = 2, ...), when a user moves inward, $nk-1 \rightarrow nk-1+1$  and  $nk \rightarrow nk-1$  with a rate of nkPk, $I\gamma$ .

The state transition rates can be obtained from the above investigation. And then, the steady-state probabilities, denoted as (S), can be obtained with the normalization condition  $\Sigma$ all S (S) = 1.

# 2.5 Performance Measures

The new call blocking probabilities for the WLAN and the cell, PB, w and PB, c can be expressed as

$$PB, W = K \sum_{k=1}^{\infty} \sum S \in \Omega \cap \Omega(S) Pw, k$$
 (11)

$$PB, C = K \sum_{k=1} \sum S \in \Omega(S) Pc, k$$
(12)

Where

$$\Omega = \{S|(S) + rw > Cw\}$$
  

$$\Omega k = \{S|(S) + rk > Cc - Cg\}.$$

The overall blocking probability PB is PB,c + PB,w. The handoff failure probability in the WLAN is

$$Pf, w = K \sum_{k=1}^{\infty} \sum_{s \in T(S)} Pw \beta k + \sum_{s \in TK} S \in TK\Pi(s) Pw \beta k + 1$$
(13)

where  $\Gamma k = \{S|Rc(S) + rk > Cc\}$ . Similarly, the handoff failure probability in the cellular network is

$$Pf = \sum_{S \in T} (S) \{ (1 - \alpha_{k+1}) + \alpha_{k+1} (1 - w_k) \} + \sum_{S \in O \cap \Gamma} \Pi(S) \alpha_{k+1} w_k$$
(14)

and 
$$P_f = P_{f,c} + P_{f,w}$$

A new call attempts to connect to the WLAN with a probability of wk when it arrives in the double-coverage area. Its attempt can fail when there are not enough resources available in the WLAN, though the call may be serviced in the cellular network if rk resources are available. The new call attempt failure probability, P'B, is

$$P_{B} = \sum_{k=1} \sum_{S \in \Omega(S)} P_{w,k} W_{k}$$

$$\tag{15}$$

A downward VHO attempt to the WLAN fails when not enough resources are available in the WLAN. It should be noted that the call requesting a VHO is not forced to be terminate if its VHO attempt fails. An upward VHO attempt to the cellular network can fails and the call terminates when there are not enough resources in the target ring to support the call, i.e., at least rk resources for the kth ring. The handoff attempt failure probability, Pf, is

$$P_{f} = \sum_{k=1}^{\infty} \sum_{s \in \Omega} (S) P_{c,k} \alpha_{k} W_{k} + \sum_{k=1}^{\infty} \sum_{s \in \Gamma_{k}} (S) P_{w} \beta_{k} + \dots$$
(16)

A call may be dropped when moving within a sector since a call moving outward requires more resources. When a user moves inward, there is no blocking because the amount of resources required by the user decreases. As shown in [2], the call dropping probability due to mobility (PD) can be kept low when the reserved resources Cg can be assigned to outward moving calls as well as handoff calls, so we do not examine PD here.

Table 1 Three cases

**IJSER** © 2013

Case	$C_c$	$C_g$	$C_w - n_w$			
1	15	2	0			
2	15	2	8			
3	15	2	2			

## 3 SIMULATION AND NUMERICAL RESULTS

As shown in Fig. 1(a), a WLAN AP is located 200 m apart from the BS and its coverage radius is 100 m. The radius of coverage of the cellular network is set to 300 m and the cell coverage area is divided into three rings.

## 3.1 Effective resources

In order to get the required resources in each ring, the same assumptions have been used as in [2]. The path loss in dB at a distance d m from the BS is  $39.95 + 43.375 \log 10d$  10. A frequency selective fading channel is considered according to a 3GPP typical case for an urban area with a receiver with 6 taps [6]. Simulation parameters are based on the 3GPP LTE system: A 5 GHz band contains 25 resource blocks (RBs), where each RB consists of 12 subcarriers with a spacing of 15 kHz. Total transmit power at the BS is 43dBm and the target BER is 0.1%. The data rate at RB *j* of user *i* is given by  $Ri = 12\log 2 (1 + \beta SNR_{i,j})$ , where 12 is the number of subcarriers in an RB,  $\beta$  = 1.5-ln(5BER) , and SNRi,j is the SNR averaged over 12 subcarriers in RB j of user i. The noise power density No is -174 dBm/Hz. The target data rate is 10 kb/s with single antenna configurations for both BS and User Equipment (UE).

When the target data rate can be supported with a partial use of an RB, the RB doesn't need to be used continuously in the time domain. The system can use time sharing in this case; the amount of resources for a user can be less than one and the remaining can be assigned to other users. Scheduling is based on the heuristic max-min method [7]. The required resources are measured by the portion of an RB required to support the target data rate, assuming an equal power distribution over the frequency band since it has been shown that water-filling has an insignificant impact after multi-user scheduling. As found in [2], rk is independent of the traffic load and cell size, and depends on the target data rate. Interestingly, if r1 is normalized to be one, R = (r1, r2, r2, r2, r3)r3) can be approximated as (1, 1.5, 2), i.e., a user in the 3rd ring may require twice the resources of a user in the innermost ring.

## 3.2 Performance Evaluation

We use  $R=(1,\,1.5,\,2)$  and set rw=1. We consider three cases for comparison as shown in Table 1. The capacity of the cellular network (Cc) is 15 and Cg is 2. The available resources for the cellular users in the WLAN, Cw=nw, is varied. The preference w3 is set to 1 so that the call requiring the most resources in the cellular network always prefers the WLAN over the cellular network. w1 and w2 are each varied in the range  $[0,\,1]$ . When there is no WLAN, PB=0.14, Pf=0.04 and their ratio, P, is 3.5. The mean call duration  $1/\mu c$  is set at 100 s.

For Case 1, P' B increases as w2 increases and converges to 1/3 which corresponds to Pw. P'f increases as well but is

kept low due to guard resources. For Case 2, there are sufficient resources in the WLAN; P'B and P'f are kept very small e.g., less than 10–3 and P'f decreases as (w2, w3) approaches (1, 1). In this case, it is best to accommodate as many cellular users as possible in the WLAN, i.e., PB and Pf are minimized when W = [111]. It should be noted that PB and Pf are always smaller when a WLAN overlaps the cell. Finally, if the available resources in the WLAN are small as in Case 3, P'B and P'f become significant and they increase as  $W \rightarrow [1\ 1\ 1]$ unlike in Case 2.

Generally, the consequences of a handoff failure are worse than those of a new call blocking. We introduce an outage probability for all connections,

$$P_{out} = P_B + 3.5 \tag{17}$$

where we use the weighting factor of 3.5 since PB is 3.5 times higher than Pf without the WLAN. Similarly, the new connection attempt failure probability, P' out, can be defined with P'B and P'f in the same way. Then, we define the overall cost as

$$Cost = P_{out} + f.....P_{out}$$
 (18)

where f is the weighting factor for the new connection attempt failure probability. Fig. 2 shows Cost versus w1 and w2 with f=3.5 for Case 3.  $W=[0\ 0.3\ 1]$  minimizes the cost for this case. The cost increases most rapidly with respect to w2, and  $W=[1\ 1\ 1]$ , i.e., unconditional preference to the WLAN by users in every ring, pays the largest cost. Thus, as the remaining resources in the WLAN decrease, it is better to for the WLAN to give higher preference to the users who requiring the most resources in the cellular network.

TABLE 2 Preference settings to minimize the cost with various values for the weighting factor  $f(w_3=1)$ .

f	0	2	4	6	8	10	12	14	16
$w_1$	0	0	0	0	0	0	0.5	0.7	0.7
$w_2$	0	0	0.4	0.7	0.9	1.0	1.0	1.0	1.0
$\frac{C_{min}}{C_{1,1}}$	0.33	0.78	0.91	0.96	0.98	0.99	1	1	1

To further investigate the impact of the weight f, Table 2 shows the optimum preference settings as well as the ratio of the minimum value of *Cost* over the value of *Cost* when W = [111]. With a small value of f, it is beneficial not to allow the cellular traffic to perform VHOs to the WLAN. Increasing f reflects greater importance of the effect of new connection attempt failures. In such a situation, more cellular traffic is allowed to attempt to move to the WLAN, i.e.,  $W \rightarrow [1 \ 1 \ 1]$ . We also note that in the considered scenario, w1 has an insignificant impact on the gain in *Cost* and can be set to 0.

## 4 Conclusions

We have proposed a new call admission control scheme for heterogeneous networks, where the users' relative preference for the WLAN changes adaptively based on the available resources in the WLAN and the location distribution of the cellular users. By modifying our ring-based model of the cell, we derived performance measures such as call blocking/ handoff failure probabilities as well as new call/handoff attempt failure probabilities. We investigated the effect of WLAN preference settings with various WLAN loading and cost weighting factors in a sector with 3 rings and a WLAN hotspot. Using the proposed CAC scheme the system cost can be greatly decreased with respect to the case where users always prefer the WLAN (i.e.,  $W = [1\ 1\ 1]$ ).

#### 5 References

- [1] D. Astély, E. Dahlman, A. Furuskär, Y. Jading, M. Lindström, and S.Parkvall, "LTE: the evolution of mobile broadband," IEEE Commun.Mag., pp. 44-51, Apr. 2009.
- [2] D. K. Kim, D. Griffith, and N. Golmie, "A novel ring-based performance analysis for call admission control in wireless networks," IEEE Commun. Lett., vol. 14, no. 4, pp. 324-326, Apr. 2010.
- [3] D. K. Kim and D. K. Sung, "Characterization of soft handoff in CDMA systems," IEEE Trans. Veh. Technol., vol. 48, no. 4, pp. 1195-1202, July 1999.
- [4] W. Song, H. Jiang, and W. Zhuang, "Performance analysis of the WLAN first scheme in cellular/WLAN interworking," IEEE Trans. Wireless Commun., vol. 6, no. 5, pp. 1932-1943, May 2007.
- [5] S. Lee, K. Kim, K. Hong, D. Griffith, Y. H. Kim, and N. Golmie, "A probabilistic call admission control algorithm for WLAN in heterogeneous wireless environment," IEEE Trans. Wireless Commun., vol. 8, no. 4, pp. 1672-1676, Apr. 2009.
- [6] 3GPP TS45.005 v8.3.0, "Radio access network; radio transmission and reception (Release 8)," Nov. 2008.
- [7] M. Tao, Y. C. Liang, and F. Zhang, "Resource allocation for delay differentiated traffic in multiuser OFDM systems," IEEE Trans. Wireless Commun., vol. 7, no. 6, pp. 2190-2201, June 2008.

# **ABOUT THE AUTHORS:**

- Samrin Shareef completed Bachelor of Technology (B.Tech) in Electronics & Communication Engineering from Jawaharlal Nehru Technological University, Hyderabad in 2008.and currently pursuing Masters of Technology (M.Tech) in Digital Electronics and Communication Systems from Jawaharlal Nehru Technological University, Kakinada.
- P.Prasanna Murali Krishna is working as Associate Professor in ECE Department and is also Head of the Department with Dr.SGIET, Markapur- Jawaharlal Nehru Technological University, Kakinada from 2001 till date.

